

Improving Physical-Layer Security for Full-duplex Radio aided Two-Way Relay Networks

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Abstract

The power allocation optimization problem is investigated for improving the physical-layer security in two-way relaying networks, where a full-duplex relay based half-jamming protocol (HJP-FDR) is considered. Specially, by introducing a power splitter factor, HJP-FDR divides the relay's power into two parts: one for forwarding the sources' signals, the other for jamming. An optimization problem for power split factor is first developed, which is proved to be concave and closed-form solution is achieved. Moreover, we formulate a power allocation problem to determine the sources' power subject to the total power constraint. Applying the achieved closed-form solutions to the above-mentioned problems, a two-stage strategy is proposed to implement the overall power allocation. Simulation results highlight the effectiveness of our proposed algorithm and indicate the necessity of optimal power allocation.

Keywords: Full-duplex, Two-way relay network, Physical-layer security, Power allocation

1. Introduction

Cooperative relays have recently been the subject of great interests among the research community owing to their abilities to achieve a broader coverage of wireless networks and obtain the diversity gain [1][2]. Depending on whether the relay broadcasts the two sources' overlapped signals, there are two types of cooperative relaying, i.e., one-way relay (OWR) [3] and two-way relay (TWR) [4]. Due to the half-duplex (HD) [5] mode, one bi-directional communication process costs four time slots for OWR. Fortunately, TWR is able to reduce the process cost to two time slots by allowing the two-way communication be performed in the same channel, simultaneously. Therefore, TWR can greatly improve the spectral efficiency [6]. Particularly, TWR based on amplify-and-forward (TWR-AF) protocol has attracted prominent attentions in many aspects due to its low-complexity in many aspects, such as relay selection [7], performance evaluation [2] [8], and resource allocation [9].

However, wireless communication systems are also open to potential eavesdroppers because of the broadcasting characteristics of wireless channels. Therefore, physical-layer security (PLS) has become one of key issues for wireless communications and attracted great attentions [10][11]. Compared with traditional higher layers' cryptographic methods, PLS technologies protect the communications against eavesdroppers through the utilization of the physical characteristics of wireless channels. In [12], a joint relay and jammer selection in TWR-AF was investigated, where two or three intermediate nodes were selected to enhance PLS. It has been shown that cooperative jamming is able to achieve considerable performance gains compared with the non-jamming schemes. Following the study, [13] proposed a Stackelberg-game-based power allocation scheme to deal with the tradeoff between the friendly jammers and the source nodes. However, the jamming process may also degrade the signal-to-noise ratio (SNR) of legitimate nodes. To overcome this problem, the authors in [14][15] have indicated that: the information leakage mainly happened in phase one and the jamming process in phase two was not necessary. Particularly, to improve the system secrecy performance, [15] adopted a jammer to interfere with the eavesdroppers in the first time slot. In this way, the impact of jamming process on legitimate links is reduced.

To further reduce the negative impact of jamming process, full-duplex radio (FDR) has recently been introduced. FD allows the node to transmit and receive signals asynchronously over the same frequency band [17]-[21]. In [16], a half-jamming protocol based on FDR (HJP-FDR) was developed. Specially, instead of adopting additional jammer, the relay was responsible for the duties of both receiving and jamming by operating in FD mode in the first time slot. The relay can perfectly remove its own jamming signal by using the methods in [18], [22] during the second time slot. Therefore, the negative impact of jamming process can be eliminated. However, only equal power allocation was presented in [16], thus its performance improvement was heavily limited. Therefore, in this paper, we develop the optimization problem to choose properly the power splitter factor and source power allocation for HJP-FDR in TWR-AF. The contributions of our work can be summarized as follows:

- **Power split factor:** An optimization problem is developed to choose properly the power splitter factor for jamming. By simplifying the objective function, we have proved the problem is concave and derived its closed-form solution.
- **Source power allocation:** We formulate a power allocation problem to determine the sources' power subject to total power constraint. By replacing the exact capacity expressions with their upper bounds, the simplification of the objective function is given. Closed-form

solution is then obtained.

- **Two-stage strategy:** It has been shown that the solution to the power splitter factor is independent with that of the source power allocation. Therefore, a two-stage strategy is finally presented to implement the joint optimization of source power and power splitter factor.

We organize the rest of our paper as follows. System model is firstly introduced in Section II. The optimal power allocation is developed in Section III. Simulation results and conclusion are given in Section IV and Section V respectively.

2. System Model

This paper considers a wireless two-way relay network with four nodes: two sources A and B , one relay R and one eavesdropper E , as shown in Fig. 1. R helps the communications of A and B , while E attempts to intercept the private information. The relay can switch between HD mode and FD mode, while both source nodes operate in HD mode. Particularly, by introducing a power splitter factor p ($0 \leq p \leq 1$), we divide the relay's power, P_R , into two parts: $(1-p)P_R$ is for forwarding the sources' signals, and pP_R is for jamming. The detailed protocol, i.e., HJP-FDR [16], is illustrated by Fig. 1.

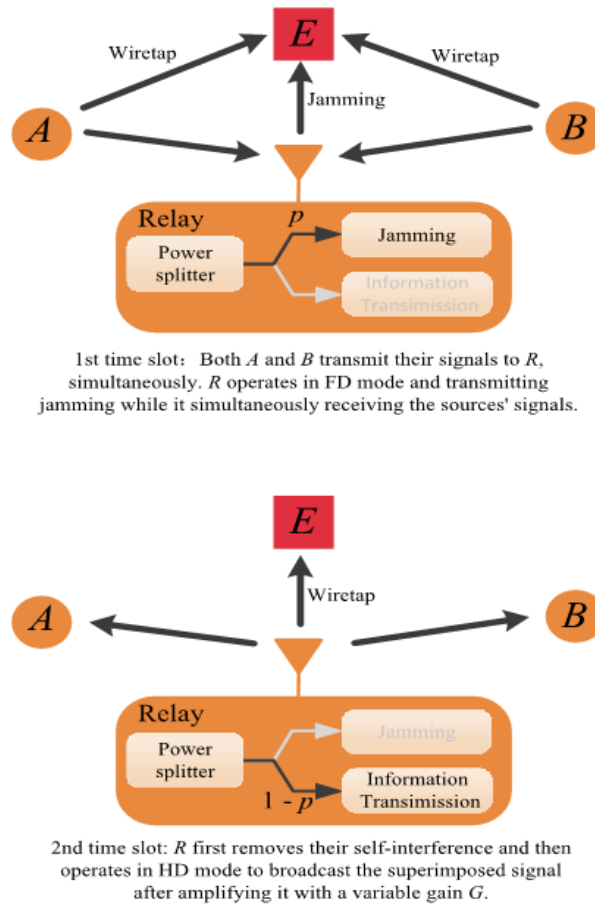


Fig. 1. System model and illustration for HJP-FDR protocol.

Obviously, HJP-FDR does not employ the jamming process in the relay's broadcasting time slot. This is because the information leakage mainly happened in phase one and the jamming process in phase two was not necessary [14] [15]. Particularly, the jamming process in HJP-FDR has no impact on the two sources, since R is able to eliminate its self-interference following the schemes in [18] [23]. Moreover, HJP-FDR does not require additional jamming node, which has greatly reduced the complexity of the whole system.

Below, we first directly present the performance of the legitimate links and then derive in detail the performance of the wiretapped links.

The instantaneous SNRs for links $B \rightarrow A$ and $A \rightarrow B$ can be given by [14]

$$\Upsilon_A^D = \frac{(1-p)P_R P_B g_{AR} g_{BR}}{((1-p)P_R + P_A)g_{AR} + P_B g_{BR}} \quad (1)$$

$$\Upsilon_B^D = \frac{(1-p)P_R P_A g_{AR} g_{BR}}{((1-p)P_R + P_B)g_{BR} + P_A g_{AR}} \quad (2)$$

where P_A , P_B denote respectively the transmit power of A , B . $g_{AR} = |h_{AR}|^2 / N_0$ and $g_{BR} = |h_{BR}|^2 / N_0$, in which h_{AR} and h_{BR} denote the channel gain of links $A-R$ and $B-R$ respectively, and N_0 is the power spectral density of the noise. Then, the capacities of the desired channels are given by

$$C_A^D = \frac{1}{2} \log_2 (1 + \Upsilon_A^D), \quad (3)$$

$$C_B^D = \frac{1}{2} \log_2 (1 + \Upsilon_B^D). \quad (4)$$

For the wiretapped links, in phase one, R operates in FD mode and receives the sources' signals with power $(1-p)P_R$, while it simultaneously broadcasts the jamming noise to E . The received signal at E in phase one can be written as

$$y_{ER}^{(1)}(t) = \sqrt{P_A} h_{AE} s_A(t) + \sqrt{P_B} h_{BE} s_B(t) + \sqrt{pP_R} h_{ER} s_J(t) + n_1(t) \quad (5)$$

where $s_A(t)$ and $s_B(t)$ denote respectively the signals transmitted by A and B with unit power. $n_1(t) \in CN(0, N_0)$ is the additive white Gaussian noise (AWGN) at E , and $s_J(t)$ is the jamming noise from R . The wiretapped SNRs from A and B in phase one are given by

$$\Upsilon_A^{E(1)} = \frac{P_B g_{BE}}{P_A g_{AE} + pP_R g_{ER} + 1}, \quad (6)$$

$$\Upsilon_B^{E(1)} = \frac{P_A g_{AE}}{P_B g_{BE} + pP_R g_{RE} + 1}. \quad (7)$$

where $g_{AE} = |h_{AE}|^2 / N_0$, $g_{BE} = |h_{BE}|^2 / N_0$, h_{AE} and h_{BE} denote respectively the channel gain of links $A-E$ and $B-E$.

In phase two, R first rebuilds and perfectly eliminates its self-interference, then forwards the signal with a variable gain by operating in HD mode. For the sake of mathematical tractability, we use a widely adopted approximation amplifying gain $G = \sqrt{(1-p)P_R / (P_A |h_{AR}|^2 + P_B |h_{BR}|^2)}$ [14], [24] by ignoring the noise statistic at the relay. It has been demonstrated that the relative error of this approximation is low even when SNR = 0 dB [25]. Thus, the received signal at E is given by

$$y_{RE}^{(2)}(t) = h_{RE} G \left(\sqrt{P_A} h_{AR} s_A(t) + \sqrt{P_B} h_{BR} s_B(t) + n_2(t) \right) + n_3(t) \quad (8)$$

where $n_2(t) \in CN(0, N_0)$ and $n_3(t) \in CN(0, N_0)$ denote respectively the AWGN at R and E . Then the wiretapped SNRs from A and B in the phase two can be written as

$$\Upsilon_A^{E(2)} = \frac{(1-p)P_R P_B g_{BR} g_{ER}}{(1-p)P_R P_A g_{ER} g_{AR} + E_{ER}} \quad (9)$$

$$\Upsilon_B^{E(2)} = \frac{(1-p)P_R P_A g_{AR} g_{ER}}{(1-p)P_R P_B g_{ER} g_{BR} + E_{ER}} \quad (10)$$

where $E_{ER} = (1-p)P_R g_{ER} + P_A g_{AR} + P_B g_{BR}$. Thus, the wiretapped rate from B and A to E can be written as

$$C_A^E = \frac{1}{2} \log_2 (1 + \Upsilon_A^E), \quad (11)$$

$$C_B^E = \frac{1}{2} \log_2 (1 + \Upsilon_B^E), \quad (12)$$

where $\Upsilon_i^E = \Upsilon_i^{E(1)} + \Upsilon_i^{E(2)}$ with $i \in \{A, B\}$.

In summary, the secrecy rates for links $B \rightarrow A$ and $A \rightarrow B$ can be presented as

$$C_A^S = C_A^D - C_A^E, \quad (13)$$

$$C_B^S = C_B^D - C_B^E. \quad (14)$$

Then, the sum secrecy rate (SSR) can be presented as

$$C^{S-sum} = C_A^S + C_B^S. \quad (15)$$

3. Power Allocation

3.1 Power Splitter Factor p

The objective of this subsection is to choose properly the power splitter factor p . Specifically, we consider the following problem:

P1: Optimal problem for p

$$\begin{aligned} \max_p C^{S-sum} \\ s.t. \quad 0 \leq p \leq 1 \end{aligned} \quad (16)$$

Since it is difficult to solve the problem in (16) straightforwardly, we suggest a suboptimal approach by approximating Υ_A^E and Υ_B^E as:

$$\Upsilon_A^E \approx \frac{P_B g_{BE}}{P_A g_{AE} + p P_R g_{ER}} + \frac{P_B g_{BR}}{P_A g_{AR}} \quad (17)$$

$$\Upsilon_B^E \approx \frac{P_A g_{AE}}{P_B g_{BE} + p P_R g_{ER}} + \frac{P_A g_{AR}}{P_B g_{BR}} \quad (18)$$

Then the problem can be represented as **P2**:

P2: Suboptimal problem for p

$$\begin{aligned} \max_p \frac{\frac{(1-p)P_R g_{AR} g_{BR}}{((1-p)P_R + P_A)g_{AR} + P_B g_{BR}}}{\frac{g_{BE}}{P_A g_{AE} + p P_R g_{ER}} + \frac{g_{BR}}{P_A g_{AR}}} \frac{\frac{(1-p)P_R g_{AR} g_{BR}}{((1-p)P_R + P_B)g_{BR} + P_A g_{AR}}}{\frac{g_{AE}}{P_B g_{BE} + p P_R g_{ER}} + \frac{g_{AR}}{P_B g_{BR}}} \\ s.t. \quad 0 \leq p \leq 1 \end{aligned} \quad (19)$$

The objective function of (19) is concave when $0 \leq p \leq 1$. Therefore, there exists a unique global optimal solution for (19). The brief proof is given as follows.

Proof: For simplicity, we use some general symbols, such as a_1 - a_5 , b_1 - b_5 , to represent the unrelated component data in (19). Thus, the objective function of (19) can be simply denoted as

$$\varphi_1(p) \times \varphi_2(p) = \frac{\frac{a_1}{1+a_2(1-p)}}{\frac{1}{a_3+a_4p}+a_5} \times \frac{\frac{b_1}{1+b_2(1-p)}}{\frac{1}{b_3+b_4p}+b_5} \quad (20)$$

We then can get what we wanted by analyzing Equ. (20). The second derivative of $\varphi_1(p)$ can be given as

$$\begin{aligned} \varphi_1''(p) = & -\frac{2a_1}{(1+a_5(a_4p+a_3))^3(1-p+a_2)^3} \times (a_2(a_4p+a_3)^3a_5^2 + (a_2(a_2+1)a_4+a_3) \\ & + (3a_2(a_2/3+p+1/3)a_3a_4 + 2a_2a_3^2 + (a_2^2 + \xi(p)a_2 + (1-p)^3)a_4^2)a_5) \end{aligned} \quad (21)$$

where $\xi(p)=3p^2-3p+2$ is a quadratic polynomial with one variable and it can be easily proved that $\xi(p) \geq \frac{4}{11} > 0$. Therefore, $\varphi_1''(p) < 0$ when $0 \leq p \leq 1$. Similarly, we have $\varphi_2''(p) < 0$ when $0 \leq p \leq 1$. As a result, the objective function of (19) is a concave function. ■

By setting the derivative of the objective function to be zero, we can obtain the optimal solution as follows:

$$\frac{1}{g_{ER}P_R} \text{Root}(\omega_1p^4 + \omega_2p^3 + \omega_3p^2 + \omega_4p + \omega_5) \quad (22)$$

where $\omega_1, \omega_2, \dots, \omega_5$ are given in (23). Obviously, (22) is a 4th-order polynomial which has only one root within $[0, 1]$. The roots of general polynomials of up to the 4th-order have closed-form expressions, e.g., Ferrari's solution and Euler's solution [26].

$$\begin{cases} \omega_1 = 2g_{AR}g_{BR} \\ \omega_2 = 3P_Ag_{AR}^2g_{BE} + 4g_{BR}(P_Ag_{AE} + P_Bg_{BE})g_{AR} + 3P_Bg_{BR}^2g_{AE} \\ \omega_3 = (2P_A^2g_{AR}g_{BR} + 6P_AP_Bg_{BR}^2)g_{AE}^2 + (6P_AP_Bg_{BE}^2 - P_AP_Rg_{BE}g_{ER})g_{AR}^2 + 2P_B^2g_{AR}g_{BE}^2g_{BR} \\ \quad + (2P_A^2g_{AR}^2g_{BE} + 12P_AP_Bg_{AR}g_{BE}g_{BR} + 2P_B^2g_{BE}g_{BR}^2 - P_BP_Rg_{BR}^2g_{ER})g_{AE} \\ \omega_4 = P_BP_A(3P_Bg_{AR}^2g_{BE}^2 + ((4P_Ag_{AR}^2 + 7P_Bg_{AR}g_{BR})g_{AE} - 2P_Rg_{AR}^2g_{ER})g_{BE}^2 \\ \quad + ((7P_Ag_{AR} + 4P_Bg_{BR})g_{AE} - 2P_Rg_{AR}g_{ER})g_{AE}g_{BR}g_{BE} + g_{AE}^2g_{BR}^2(3P_Ag_{AE} - 2P_Rg_{ER})) \\ \omega_5 = 2P_B(P_Ag_{BE}(P_Bg_{BE} - P_Rg_{ER}/2)g_{AE}^2 + P_AP_Bg_{AR}g_{AE}g_{BE}^2 - P_BP_Rg_{AR}g_{ER}g_{BE}^2/2) \\ \quad \cdot (g_{AR}g_{BE} + g_{BR}g_{AE})P_A \end{cases} \quad (23)$$

3.2 Source Power Allocation

In this subsection, a power allocation optimization problem is developed for maximizing SSR subject to total power constraint P_T . Specifically, we consider the following problem:

P3: Optimal problem for source power

$$\begin{aligned} \max_{P_A, P_B, P_R} & C^{S-sum} \\ \text{s.t.} & P_A, P_B, P_R \geq 0 \\ & P_A + P_B + P_R \leq P_T \end{aligned} \quad (24)$$

Similarly, it is hard to achieve the optimal solution to the problem in (24), directly. We consider the simplification of the objective function by replacing the exact capacity expressions with their upper bounds as:

$$\Upsilon_{AR}^D < (1-p)P_Rg_{AR}, \quad (25)$$

$$\Upsilon_{BR}^D < (1-p)P_R g_{BR}, \quad (26)$$

$$\Upsilon_{AR}^E < \frac{P_B g_{BR}}{P_A g_{AR}}, \quad (27)$$

$$\Upsilon_{BR}^E < \frac{P_A g_{AR}}{P_B g_{BR}}. \quad (28)$$

Through large amounts of numerical simulations, we have found that the value of P_R is around $P_T/2$ when the problem (24) approach its optimal value. Therefore, to further simplify the problem, we set $P_R = P_T/2$. Then, **P3** can be written as **P4**.

P4: Suboptimal problem for source power

$$\begin{aligned} \max_{P_A, P_B, P_R} \quad & \frac{1+(1-p)P_T g_{AR}/2}{1+P_B g_{BR}/P_A g_{AR}} \cdot \frac{1+(1-p)P_T g_{BR}/2}{1+P_A g_{AR}/P_B g_{BR}} \\ \text{s.t.} \quad & P_A, P_B, P_R \geq 0 \\ & P_A + P_B + P_R \leq P_T \end{aligned} \quad (29)$$

The above problem is non-convex due to the non-convex objective function. However, we still will be able to find the solution to this problem by resorting to the KKT conditions. The Lagrangian function of **P4** is presented as follows:

$$\begin{aligned} \zeta(P_A, P_B, \lambda_1, \lambda_2, \lambda_3) \\ = \frac{1+(1-p)P_T g_{AR}/2}{1+P_B g_{BR}/P_A g_{AR}} \cdot \frac{1+(1-p)P_T g_{BR}/2}{1+P_A g_{AR}/P_B g_{BR}} - \lambda_1 P_A - \lambda_2 P_B + \lambda_3 (P_A + P_B - P_T/2) \end{aligned} \quad (30)$$

The solution to **P4** must satisfy the KKT conditions [27] as

$$\begin{cases} \frac{\partial \zeta}{\partial P_A} = \frac{g_{BR} g_{AR} (2+P_T(1-p)g_{BR})(2-P_T(1-p)g_{AR})(P_A g_{AR} - P_B g_{BR})P_B}{4(P_A g_{AR} + P_B g_{BR})^3} - \lambda_1 - \lambda_3 = 0 \\ \frac{\partial \zeta}{\partial P_B} = \frac{g_{BR} g_{AR} (2+P_T(1-p)g_{BR})(2+P_T(1-p)g_{AR})(P_A g_{AR} - P_B g_{BR})P_A}{4(P_A g_{AR} + P_B g_{BR})^3} - \lambda_2 - \lambda_3 = 0 \\ P_A + P_B \leq P_T/2, \lambda_3 (P_A + P_B - P_T/2) = 0 \\ \lambda_1 P_A = 0, \lambda_2 P_B = 0, P_A, P_B, \lambda_1, \lambda_2, \lambda_3 \geq 0 \end{cases} \quad (31)$$

Then, a unique optimal solution to **P4** is achieved by solving the KKT conditions in (31):

$$P_A^* = \frac{1}{2} \frac{g_{BR}}{g_{AR} + g_{BR}} P_T, \quad (32)$$

$$P_B^* = \frac{1}{2} \frac{g_{AR}}{g_{AR} + g_{BR}} P_T. \quad (33)$$

From (33), we can find that the source power allocation is independent from p . Therefore, the overall power allocation can be simply implemented by the following two-stage strategy.

- Step I: Allocate the power of the sources and relay according to (33).
- Step II: Based on the power levels of each nodes, figure out the power splitter factor p according to (22).

It is worth noting that our suboptimal design is very close to the optimal counterpart as later shown in simulations, although some approximations are adopted.

3.3 Power Allocation without Eavesdropper's CSI

In this subsection, we consider a more realistic case where the CSI of the eavesdropper is not available. From (32) and (33), we can observe that the source power allocation is independent from CSI of E . Therefore, the proposed source power allocation can be adopted directly. However, the power splitter factor, p , determined by the CSI of E , as shown in (22), which is invalid in this case. Thus, for simplicity, we set $p=0.25$ (as later shown in Fig. 7, $p \leq 0.25$ in most cases.). Finally, the power allocation for the case without eavesdropper's CSI can be given by:

$$\begin{cases} P_A^* = \frac{1}{2} \frac{g_{BR}}{g_{AR} + g_{BR}} P_T \\ P_B^* = \frac{1}{2} \frac{g_{AR}}{g_{AR} + g_{BR}} P_T \\ P_R^* = P_T / 2 \\ \rho = 0.25 \end{cases} \quad (34)$$

4. Numerical Results

Computer simulations are performed to examine the performance of the proposed power allocation based on the simulation model as shown in Fig. 2. Specifically, the coordinate of node i is denoted by (x_i, y_i) , where x_i denotes horizontal coordinate and y_i denotes vertical coordinate. Set $y_R=0.5$, $(x_A, y_A) = (-0.5, 0)$ and $(x_B, y_B) = (0.5, 0)$. We consider Rician fading with parameter $K = 3$. Let the path loss coefficient $\alpha = 3$. For the equal power allocation, $P_A = P_B = P_R = P_T/3$, $p=0.5$.

SSR is presented in Fig. 3 against x_E under the equal power allocation (EPA) [16] and the proposed power allocation with $P_T=20\text{dBm}$, $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-0.5 \sim 0.5, 0)$. Moreover, to further examine the effectiveness of the proposed algorithm, the optimal SSR under numerical calculation is also presented. The optimal (Numerical) result here is obtained by using the computer to search the optimal set of $\{p, P_R, P_A, P_B\}$ from the entire feasible region. From Fig. 3, we can observe that the proposed power allocation has similar performance to the optimal counterpart and outperforms EPA. In this scenario, we can observe that SSR falls obviously when E is close to the sources.

Moreover, by setting $P_T=20\text{dBm}$, $(x_R, y_R) = (-0.5 \sim 0.5, 0.5)$ and $(x_E, y_E) = (0, 0)$, we compare SSR of the proposed power allocation with EPA against x_R in Fig. 4. Similarly, the optimal SSR under numerical calculation is also presented. As shown in Fig. 4, similar to Fig. 3, the proposed power allocation has similar performance as the optimal counterpart and is well ahead of EPA. For EPA, SSR decreases rapidly when R locates closely to the sources. While for the proposed power allocation, SSR is still showing a high level, indicating the necessity of optimal power allocation.

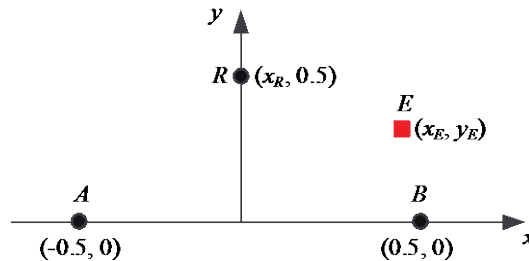


Fig. 2. Simulation model.

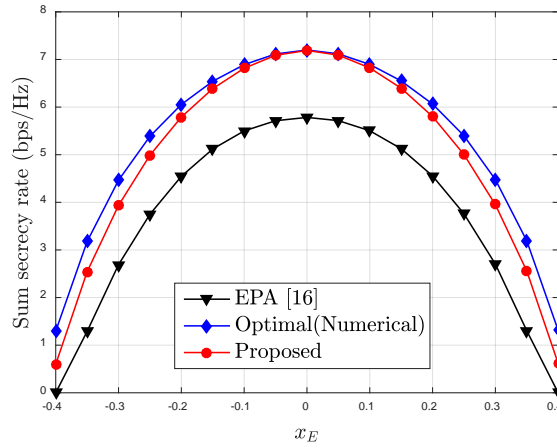


Fig. 3. SSR against d_{AE} with $P_T=20\text{dBm}$, $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$, $(x_E, y_E) = (-0.5 \sim 0.5, 0)$.

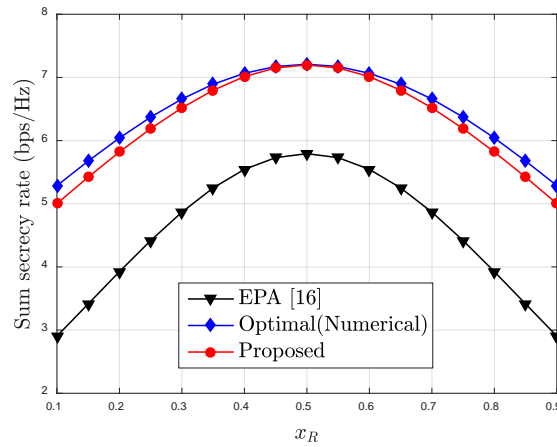


Fig. 4. SSR against d_{AR} with $P_T=20\text{dBm}$, $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (-0.5 \sim 0.5, 0.5)$, $(x_E, y_E) = (0, 0)$.

Set $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-1 \sim 1, 0 \sim 2)$, that is, E traverses the two-dimensional area. SSRs of EPA and the proposed power allocation are plotted in **Fig. 5** and **Fig. 6** against the location of E , where the colors stand for different values of SSR. As can be seen from the two figures, the proposed power allocation has brought considerable performance gain compared with EPA for any locations E , highlighting the effectiveness of our design.

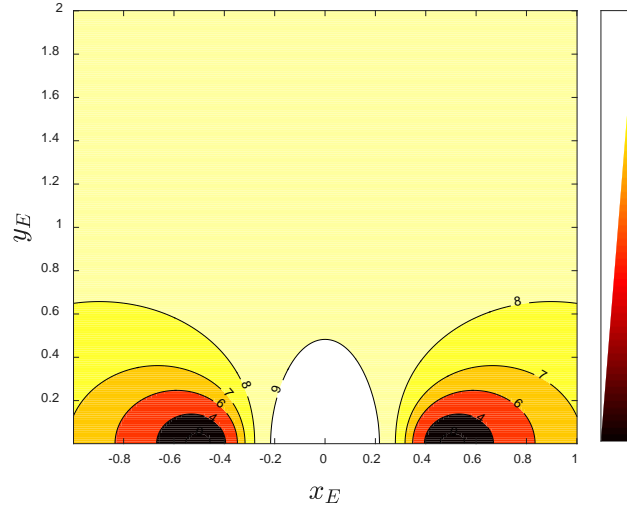


Fig. 5. SSR of EPA [16] against the location of E with $P_T=20\text{dBm}$, $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-1\sim 1, 0\sim 2)$ (different colors stand for different values of SSR).

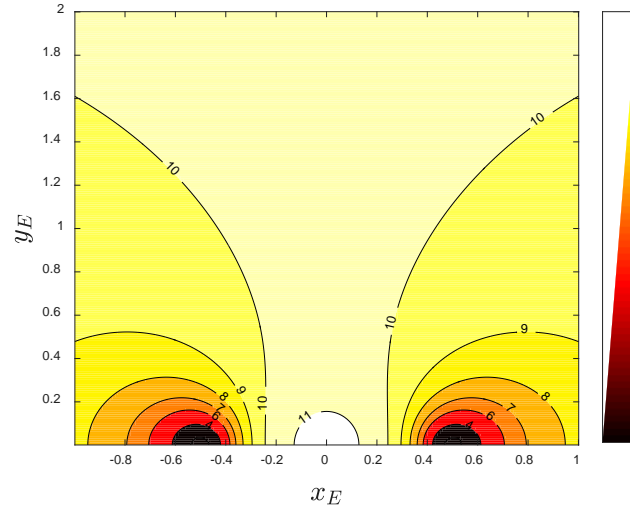


Fig. 6. SSR of the proposed power allocation against the location of E with $P_T=20\text{dBm}$, $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-1\sim 1, 0\sim 2)$.

To present the impact of the location of E on the power of jamming, **Fig. 7** plots the power splitter factor under the proposed power allocation against the location of E by using the same simulation parameter settings as **Fig. 5**. In **Fig. 7**, the colors stand for different values of p . As can be observed from **Fig. 7**, p increases, i.e., more power is allocated for jamming, when E locates closely to the sources. And the jamming power decreases when E locates in the middle between A and B , or close to the relay.

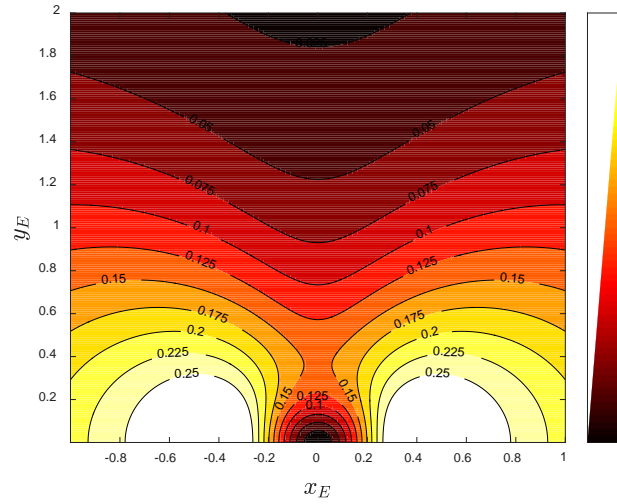


Fig. 7. Power splitter factor under the proposed power allocation against the location of E with $P_T=20\text{dBm}$, $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-1\sim 1, 0\sim 2)$.

To present the impact of P_T on SSR, **Fig. 8** plots SSR against P_T with $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$, $(x_E, y_E) = (0, 0)$. As shown in **Fig. 8**, the proposed power allocation is able to achieve non-negative when $P_T > 7.5\text{dBm}$, while EPA requires $P_T > 9.5\text{dBm}$. Secrecy performance is improved by 2dB with the increase of P_T .

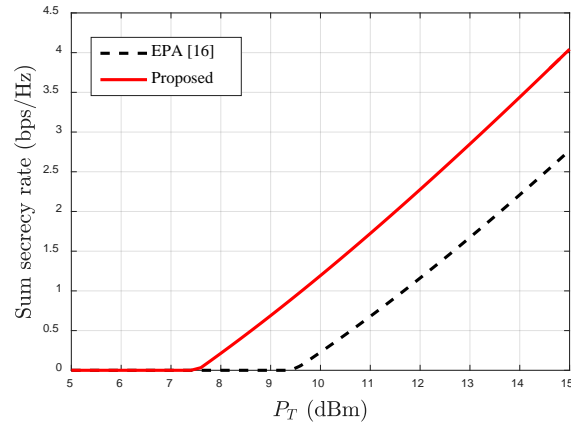


Fig. 8. SSR against P_T with $(x_A, y_A) = (-0.5, 0)$, $(x_B, y_B) = (0.5, 0)$, $(x_R, y_R) = (0, 0.5)$, $(x_E, y_E) = (0, 0)$.

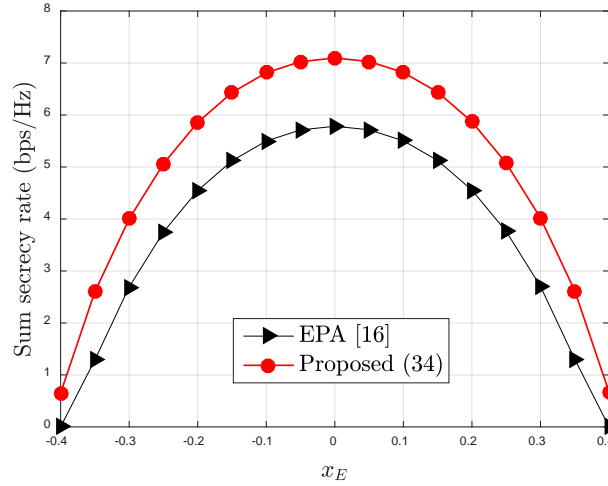


Fig. 9. SSR of EPA and the proposed power allocation in (34) without CSI of E .

Finally, we consider the case where CSI of E is not available. The SSR is presented in **Fig. 9** against x_E under EPA and the proposed power allocation in (34) with $P_T=20\text{dBm}$, $(x_R, y_R) = (0, 0.5)$ and $(x_E, y_E) = (-0.5 \sim 0.5, 0)$. From **Fig. 9**, we can observe that the proposed power allocation still brings considerable performance gain compared with EPA in [16], even though CSI of E is not available.

5. Conclusion

A two-stage power allocation strategy for HJP-FDR in TWR-AF is proposed to implement the joint optimization of source power and power splitter factor by using the fact that the solution to the power splitter factor is independent with that to source power allocation, which is derived by developing and solving the optimization problems for power split factor and source power allocation. Moreover, the case without CSI of E has also been taken into consideration. From simulations, we have shown that our proposed algorithm is very close to the optimal counterpart and outperforms equal power allocation, highlighting the necessity of optimal power allocation. It is important to note that the proposed scheme in this paper can also be applied to some specified networks, such as satellite communications where the satellite acts as the relay, multiple relay networks [28], cognitive radio networks [29], and 5G networks [30] [31]. The specific issue of signal synchronization should be addressed when using this method for satellite communications applications.

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